

Contents

Chapter 0 – Prerequisites2

Chapter 1 – Preparing the Environment3

Chapter 2 – Writing a Dummy Popup5

Chapter 3 – Writing a “LOL” Popup10

Chapter 4 – Using the Stack for a Popup13

Chapter 5 – The Easy Way – Part 117

Chapter 6 – The Easy Way – Part 220

Ending Words & References.....28

Chapter 0

Prerequisites

Tools

- Arwin (<http://www.vividmachines.com/shellcode/arwin.c>)
- Code::Blocks (<http://www.codeblocks.org/>)
- OllyDbg (<http://www.ollydbg.de/version2.html>)

Goals

- Create minimalistic shellcode
- Include custom text input

This is for educational purposes only.



Chapter 1

Preparing the Environment

First download and install the Code Blocks IDE environment, which you're going to use to compile the Arwin source code. Simply download "arwin.c" from the link in the previous chapter and open this in Code Blocks or your own preferred C compiler.

When you've opened this file in Code Blocks, browse to the "Build" menu and choose "Build".

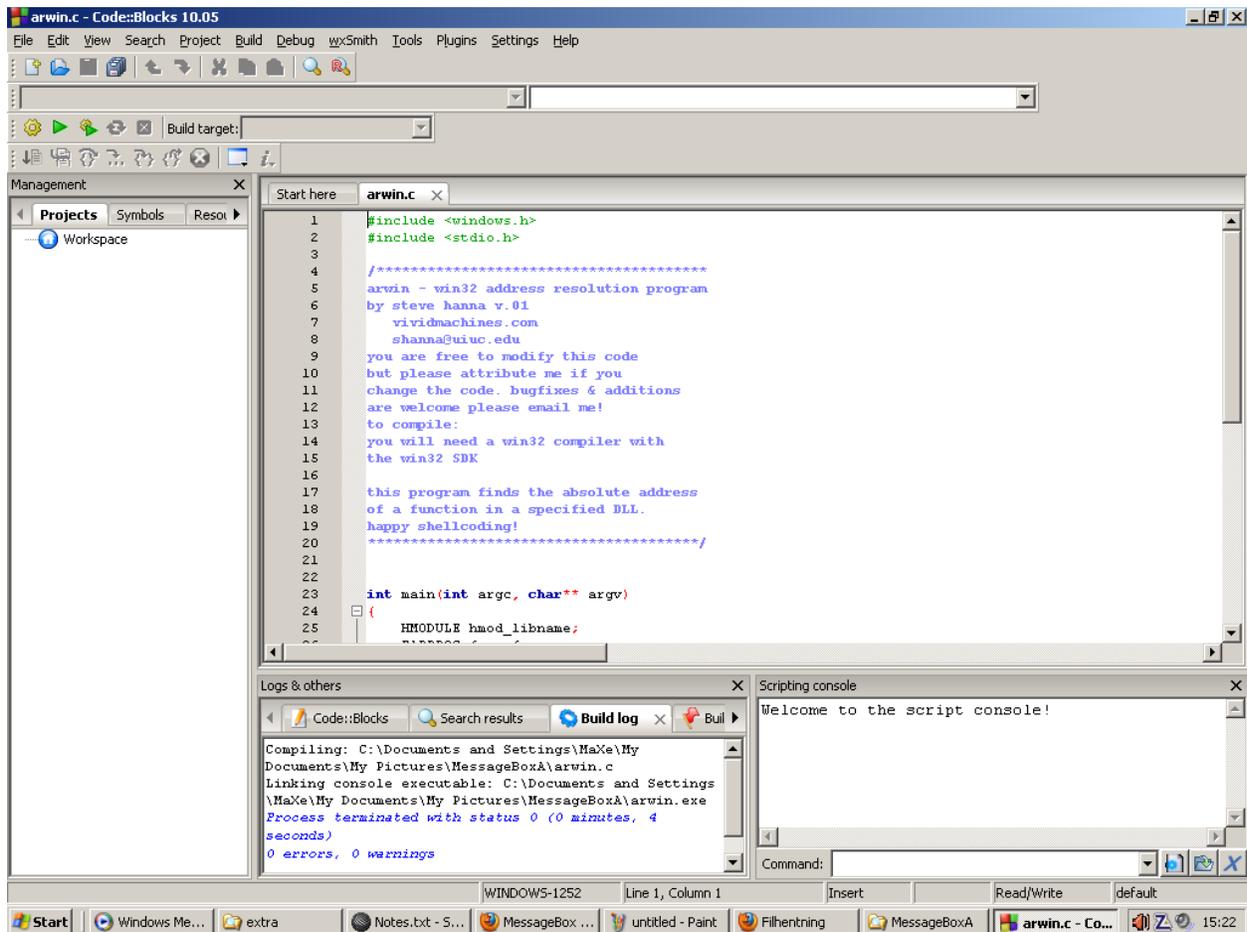
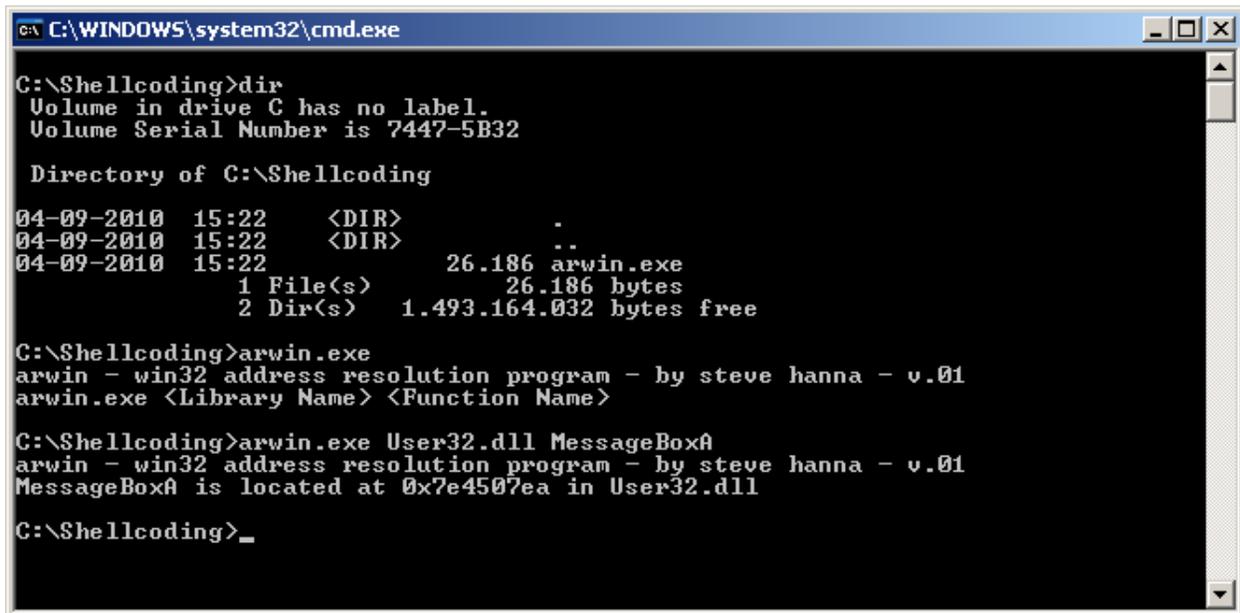


Figure 1.0.1 – Compiling Arwin.c in Code::Blocks

Now you're ready to use Arwin, which will be located in the same directory as where the source code is located. Browse to this location in a command prompt window. If the directory is located deep within your computer, consider copying arwin.exe to a directory near the root.

Then issue the following command without quotes: "arwin.exe User32.dll MessageBoxA".

When you've done this, you'll be searching for the MessageBoxA function within User32.dll



```
C:\WINDOWS\system32\cmd.exe
C:\Shellcoding>dir
Volume in drive C has no label.
Volume Serial Number is 7447-5B32

Directory of C:\Shellcoding
04-09-2010  15:22    <DIR>          .
04-09-2010  15:22    <DIR>          ..
04-09-2010  15:22                26,186 arwin.exe
               1 File(s)                26,186 bytes
               2 Dir(s)      1,493,164,032 bytes free

C:\Shellcoding>arwin.exe
arwin - win32 address resolution program - by steve hanna - v.01
arwin.exe <Library Name> <Function Name>

C:\Shellcoding>arwin.exe User32.dll MessageBoxA
arwin - win32 address resolution program - by steve hanna - v.01
MessageBoxA is located at 0x7e4507ea in User32.dll

C:\Shellcoding>_
```

Figure 1.0.2 – Using Arwin.exe to find the MessageBoxA memory address

If you've done this right, you'll see a text string like: MessageBoxA is located at ...

It is important to note, that this address is not static and it is therefore usually not the same on another machine. If you're running Windows Vista or 7, then the address will change each time you reboot your computer. If you're running certain IPS's on Windows XP, this will occur too.

After remembering these precautions, note down the memory address Arwin returned.

In our shellcode we're going to assume User32.dll is already loaded within the given executable which we are injecting our own shellcode into directly and manually. A file which already has this DLL file loaded when it is run is: GenuineCheck.exe

You can of course, use other files as well but I will be using this executable as an example.

(Any executable file with a graphical user interface should have User32.dll loaded.)

Chapter 2

Writing a Dummy Popup

Open OllyDbg and if you haven't downloaded this program yet, then download it and make sure you get version 2.0 which has a lot of improvements and new features. One of the new features is the ability to change the PE header directly, which we're going to do later on.

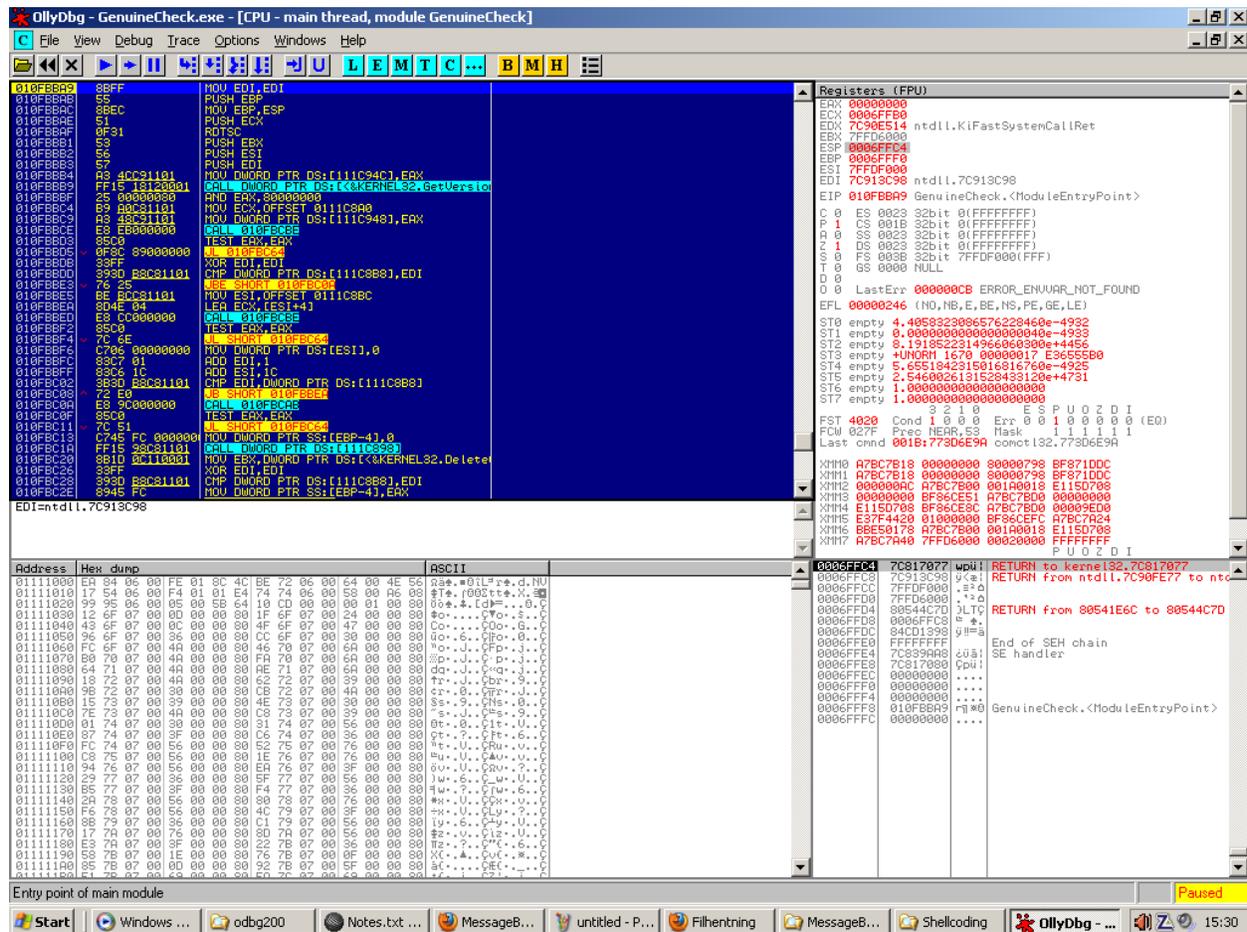


Figure 2.0.1 – An executable file within OllyDbg

Simply open your executable file as shown in Figure 2.0.1, and take a quick look at it. Don't try to understand what all of it means. In the lower right we have our stack, which stores values, variables, arguments, pointers, etc.

Essentially the same, hexadecimal values because binary values are harder to read.

In the lower left we have our "Hex Dump" window, which is useful when we're following perhaps a dynamic change of our shellcode, such as an encoding sub-routine.

The upper right is our registers and flags, while the upper left is the "disassembly window" where we can single-step aka trace through our assembly opcodes, aka (CPU) instructions.

Basically when we press F7 once, then one instruction is executed by your computer.

This may add or remove values from the stack, alter the registers, and of course if more opcodes aka instructions are used, then the hardware of the computer may be used along with perhaps API libraries to show pop boxes, connect to a remote attacker, or shut the computer down.

In essence the only limitation with Assembly is the hardware and your imagination.

Since we're not going to use the original program behind, change a good amount of the first couple of instructions visible to NOP's, so it is easier to see your own injected shellcode.

1. Select aka mark a region of opcodes shown in the "disassembly window".
2. Right click and browse to "Edit" then select: "Fill with NOPs".

All you should see is a lot of "NOP" instructions now. Press F7 a couple of times and see how you execute one NOP at a time. This opcode doesn't do anything, so don't worry if you execute many of them. As you may see, EIP changes each time you execute one NOP. This register that you can see on the right, points to the current instruction you're going to execute next.

F8 is used to jump over "CALL" opcodes but also single-step through our shellcode as well, while F9 executes the entire code until an event occurs, such as a popup box requiring attention.

Our initial dummy popup code looks like this:

```
MOV EAX, 0x7e4507ea
XOR EBX, EBX
PUSH EBX
PUSH EBX
PUSH EBX
PUSH EBX
CALL EAX
```

Figure 2.0.2 – Dummy Popup Shellcode

As you can see, I'm using this memory address: 0x7e4507ea to call the MessageBoxA function. Make sure not to use this address since it will most likely be different on your system. Use the address you found with the Arwin tool mentioned earlier.

The reason why we begin with a “dummy” is to make it easier to understand, how the call and assigned variables function together, inside your computer. Here is an explanation of the code:

- 1) Change EAX to 0x7e4507ea
- 2) XOR EBX with EBX. This alters EBX to 0. (zero)
- 3) Push the value of EBX to the stack.
- 4) Push the value of EBX to the stack.
- 5) Push the value of EBX to the stack.
- 6) Push the value of EBX to the stack.
- 7) Call the memory address which EAX is pointing to.

Note: In our case, we're calling MessageBoxA()

Why push the value of EBX 4 times to the stack? Because MessageBoxA takes 4 arguments!

Here's the MSDN syntax:

```
int WINAPI MessageBox(  
    __in_opt  HWND hwnd,  
    __in_opt  LPCTSTR lpText,  
    __in_opt  LPCTSTR lpCaption,  
    __in      UINT uType  
);
```

Figure 2.0.3 – MessageBoxA syntax

Here's the full resource: <http://msdn.microsoft.com/en-us/library/ms645505%28VS.85%29.aspx>

I will explain what the arguments means later on and how to use them with Assembly code. But for now you should double-click a NOP instruction after or where the current Instruction Pointer (EIP) is pointing to. If you don't know where EIP is pointing to, look at the very left side in your debugger where the memory addresses are.

The memory address which is highlighted is the next instruction which will be executed when you hit and press F7. When you've double-clicked a NOP instruction, begin to write the code in yourself. Make sure to use your own memory address for MessageBoxA as described earlier.

When you've done this, press F7 once and see how EAX now contains your memory address which you used Arwin to find, if you press F7 again then EBX will become 0 and if you press F7 4 times more the stack will have 4 new values of 0 (zero) each.

Now you're at the “CALL EAX” instruction. If you want to, you can press F7 all the way through the function call, but if you just need to see your message box work or not, then press F8 and watch the beautiful yet very empty popup box.

If your popup looks like the one included in the screenshot below, then you've done right.

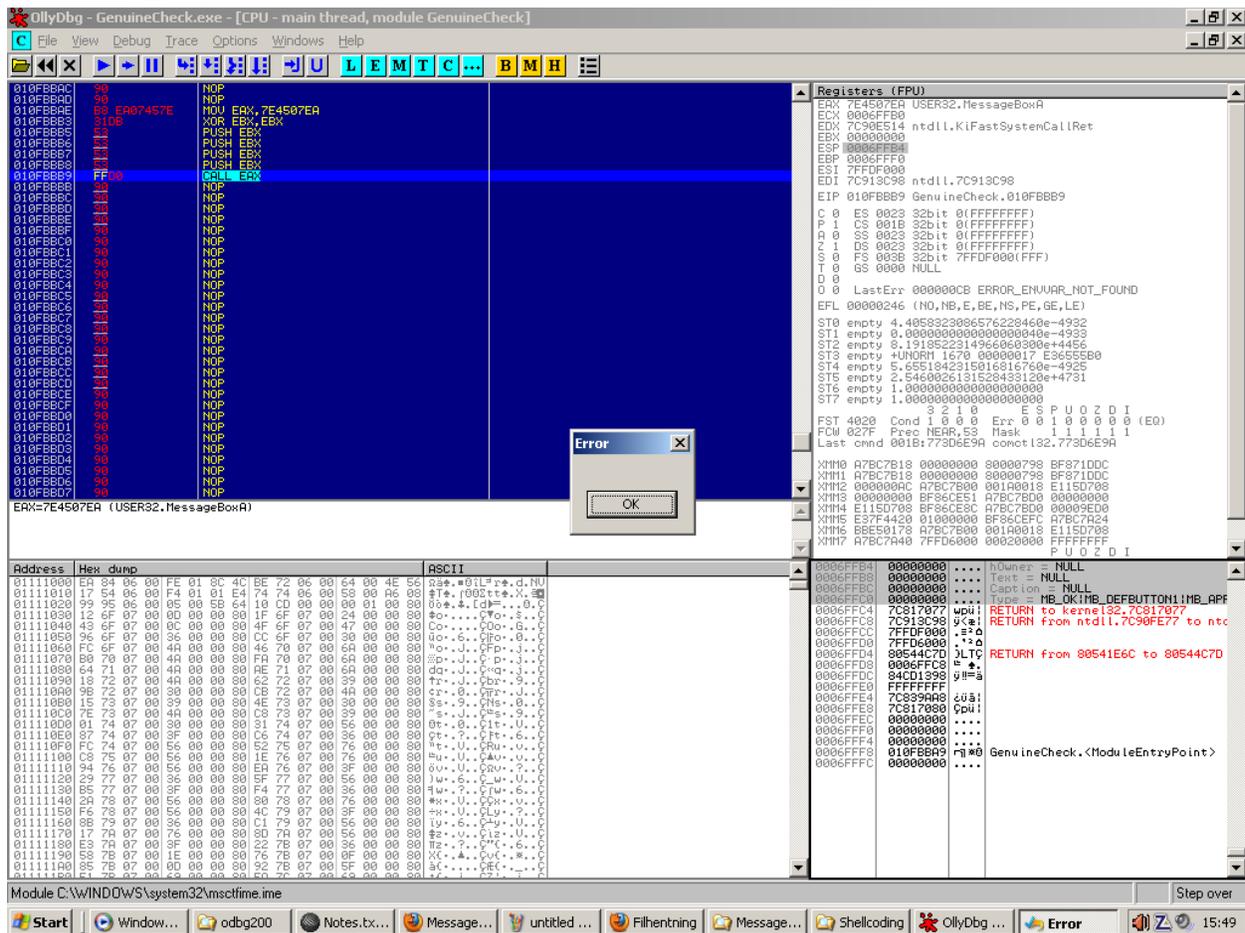


Figure 2.0.4 – A Dummy Popup working in OllyDbg

This is of course, not very impressive. Even a monkey can be taught how to do this, so we'll continue onwards and learn how to provide a custom text message now, which there are two ways which you can use in order to do this.

Before we do this, mark the shellcode you wrote. Then right-click it, browse to “Edit” and select “Binary copy” in order to save it. Paste the output into notepad and remove the spaces between.

This is the (hexadecimal) binary version of your recently created popup box. If you re-open the executable file you used by pressing CTRL+F2, overwrite the first couple of instructions with NOPS and then select a good amount of NOPS to overwrite with your binary code, then you can enter the “Edit” menu by right-clicking again and now you select “Binary paste”.

Whenever you use “Binary paste”, make sure to select **enough** opcodes!

If you don't, then only a part of your shellcode will be shown or some of the opcodes may be incomplete and therefore they may resemble something else which it shouldn't. Therefore you should as previously mentioned, always select enough opcodes aka instructions.

Before we continue I'd like to explain more about our dummy popup shellcode.

Q: Why is XOR used?

A: Because this avoids 0-bytes (00), which otherwise kills shellcode.

Q: What is the Stack used for in this scenario?

A: It's used to store arguments for MessageBoxA and possibly even custom text!

Q: I pressed F7 at "CALL EAX" and noticed a value was pushed onto the stack, why?

A: This value is the current EIP (Instruction Pointer) which is used to return back to the original code when a "RET" opcode is executed. If this value wasn't pushed onto the stack, then the computer would return to whatever other value is stored on the stack, resulting in massive failure.

Q: I saw something like this on the stack, what does it mean?

0006FFB4	00000000	; hOwner = NULL
0006FFB8	00000000	; Text = NULL
0006FFBC	00000000	; Caption = NULL
0006FFC0	00000000	; Type = MB_OK MB_DEFBUTTON1 MB_APPLMODAL

A: hOwner is the owner of the window, we'll use zero since we can. Text is a pointer to ASCII text stored somewhere in the memory, caption is the same as text and type means what kind of buttons (and more) that the popup box should display, do, require, etc.

An "ASCII Pointer" generally points to a memory address, containing a string in ASCII format with a 00 byte in the end to close the string. All this functionality will be explained in the next couple of chapters.

You might want to try this out yourself, so copy the code below and adapt it to your needs.

Address	Hex dump	Command
010FBBB1	B8 EA07457E	MOV EAX,7E4507EA
010FBBB6	31F6	XOR ESI,ESI
010FBBB8	31FF	XOR EDI,EDI
010FBBBA	BB 5D5E5D11	MOV EBX,115D5E5D
010FBBBF	81F3 11111111	XOR EBX,11111111
010FBBC5	53	PUSH EBX
010FBBC6	89E7	MOV EDI,ESP
010FBBC8	56	PUSH ESI
010FBBC9	57	PUSH EDI
010FBBCA	57	PUSH EDI
010FBBCB	56	PUSH ESI
010FBBCD	FFD0	CALL EAX

Figure 3.0.3 – Custom “LOL” Popup Code

I should mention that you actually don't need the XOR EDI, EDI opcode but I added it to make things a bit more clear. The only thing you need to change in the above code is MOV EAX to include the memory address for MessageBoxA which you found earlier using Arwin.

Do not try to copy these opcodes directly into OllyDbg, write them yourself. Pay attention to the opcode with our custom string, because if you write it wrong then our string may be displayed wrong or not at all. It's very important you do things right, since there's no margin for error.

When you've created your “LOL” popup box, do a binary copy which you can use later on:

B8 EA 07 45 7E 31 F6 31 FF BB 5D 5E 5D 11 81 F3 11 11 11 11 53 89 E7 56 57 57 56 FF D0

Figure 3.0.4 – Binary opcodes for a “LOL” popup

As you can see for yourself, we've used 29 bytes to create a message box saying “LOL”. It does not get any more efficient than this so make sure you understand it before continuing onto the next chapter which will include more advanced shellcode.

Chapter 4

Using the Stack for a Popup

After creating our popup with the words “LOL”, we’re ready to progress onto something a bit more advanced. The text string I’m going to use is: “Hello, this is MaXe from InterN0T.net”.

In short we’re going to push the string to the stack and then point to it. Easier said than done, but it is far from impossible. All it takes is time, and you got plenty of that to learn this. You might wonder why learn how to call a popup box? Well if you can do that, then you’re able to learn by yourself how to use the rest of the API!

First we have to encode our string in hexadecimal, if you use The XSSOR, make sure to use “Machine Hex Encoding” and remove these “\x” so the string looks similar to the example.

48 65 6c 6c 6f 2c 20 74 68 69 73 20 69 73 20 4d 61 58 65 20 66 72 6f 6d 20 49 6e 74 65 72 4e 30 54 2e 6e 65 74

Figure 4.0.1 – Hexadecimal Encoding of our String

The above characters could be opcodes but they are in fact our string in its encoded form. Since we’re going to push the string to the stack we need to do more work with it, yes this is painful but in the end it’ll work and show the string as it should.

To begin with we add a 00-byte in the end of our string. Furthermore our string should be dividable by 4, if it isn’t we add more 00-bytes. In our case after the first 00-byte we added, we notice that we need to add 2 more bytes to make our string more workable.

So there’s 3x 00-bytes after the last byte (74), and to save ourselves trouble from doubling or tripling the work we have to do with the string due to the Little Endian Architecture, we read the 4 last bytes and put these in 1 row and then the next 4 bytes in the next row.

The first row with the 3x 00-bytes added looks like this: 00 00 00 74, the next row: 65 6E 2E 54.

Continue with this procedure until you reach the last 4 bytes and put them on the last row. The reason why we have to do this is because the last byte is read first on the IA-32 platform. Now it will probably become even more confusing later on, but if it works then you did it right.

In the screenshot on the previous page, you're able to see that the string is pushed in yes, reverse order which is exactly how it should because this is how the Intel Architecture works. I know it's seems strange, but I promise there are more obscure things out there.

In essence row 2 to 10 was just pushed to the stack, while row 1 contained a 00 byte which had to be encoded which I did by XOR'ing the string with 1111 1111. Exactly the same trick I used in our "LOL" popup code, except that we're using a longer string this time and that it's only the first 4 bytes (actually the 4 last) that needs to be encoded.

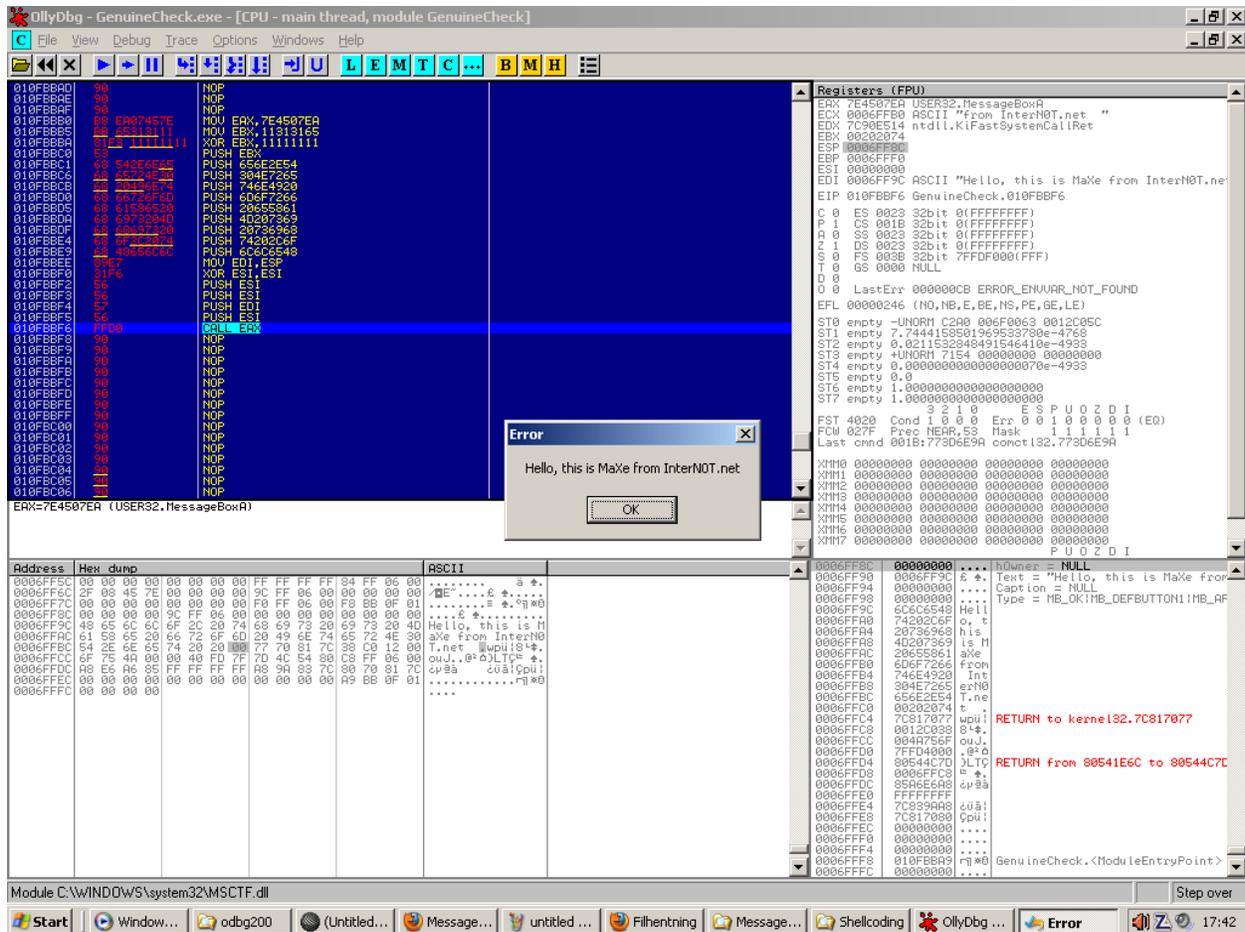


Figure 4.0.4 – Working custom popup box with no 00-bytes

When our string is pushed onto the stack, the stack pointer which points to our string is copied into the EDI register for future reference, then ESI (0) is pushed onto the stack twice. Then EDI is pushed which contains our ASCII (stack) pointer to our custom string, and finally ESI again (0) in order to complete the amount of necessary arguments.

Finally EAX is called and our popup box should be executed without any problems.

As per usual the code is available in the example below for you to use.

010FBBB0	B8 EA07457E	MOV EAX,7E4507EA
010FBBB5	BB 65313111	MOV EBX,11313165
010FBBBA	81F3 11111111	XOR EBX,11111111
010FBBC0	53	PUSH EBX
010FBBC1	68 542E6E65	PUSH 656E2E54
010FBBC6	68 65724E30	PUSH 304E7265
010FBBCB	68 20496E74	PUSH 746E4920
010FBBD0	68 66726F6D	PUSH 6D6F7266
010FBBD5	68 61586520	PUSH 20655861
010FBBD A	68 6973204D	PUSH 4D207369
010FBBDF	68 68697320	PUSH 20736968
010FBBE4	68 6F2C2074	PUSH 74202C6F
010FBBE9	68 48656C6C	PUSH 6C6C6548
010FBBEE	89E7	MOV EDI,ESP
010FBBF0	31F6	XOR ESI,ESI
010FBBF2	56	PUSH ESI
010FBBF3	56	PUSH ESI
010FBBF4	57	PUSH EDI
010FBBF5	56	PUSH ESI
010FBBF6	FFD0	CALL EAX

Figure 4.0.5 – Custom Popup Box Code

If you want to try this code right out of the box, on your own machine then simply edit the first instruction to the memory address pointing to MessageBoxA() on your system, and then enter each opcode manually into your debugger.

All of the push arguments are a bit tedious to push, so copy these one by one and insert them.

If you're really lazy, then you can use the binary codes below which are the push opcodes that pushes row 2 to 10 of our string. Remember, to select a large amount of opcodes in your debugger before your paste it in, and afterwards double-check all of the instructions indeed were pasted correctly into your debugger.

The fastest way to check this, is to check the first opcode is correct and that the last opcode is correct, and of course that the entire shellcode looks like the code in figure 4.0.5

68542E6E656865724E306820496E746866726F6D6861586520686973204D6868697320 686F2C20746848656C6C
--

Figure 4.0.6 – Binary Push Opcodes (Row 2 to 10)

Play a little with this and when you feel ready to continue onto the next chapter which is going to use another approach, then make sure you understood the previous parts because the next way is the hardest but also somewhat the most interesting.

Chapter 5

The Easy Way – Part 1

In this chapter we're going to simply add our string to the end of our shellcode. This means that we'll have a 00-byte in the end of our shellcode which may not work in real buffer overflow scenarios but in this test case it doesn't really matter. We'll solve that problem later on!

Now you may remember that we need to point to our string somehow, and that this pointer uses memory addresses. Just copy "EIP" into another register, is what you may think. This isn't possible so we actually have to calculate the EIP, no kidding.

There are many ways to Rome and my method is just one of them. It's made of pure logic and how the Assembly opcode "CALL" functions. This opcode pushes the current instruction pointer when executed and then it jumps to the memory address or register defined.

We can use this to our advantage by taking this value from the stack and storing it in a register, but we have to be careful because we can't just take a value from the stack without putting it back, since this may cause unexpected errors and more in real shellcode. (I.e. backdoors)

Take a look at this sub-routine I made: (Pseudo Code)

```
PUSH
---- SOI ---- (Start of Instructions)
POP EDI
PUSH EDI
RETN
NOP
---- EOI ---- (End of Instructions)
CALL ESP
```

Figure 5.0.1 – Calculating EIP Sub-routine

In short: Push the binary values between SOI and EOI to the stack, and then call the stack which jumps directly to the stack, take the EIP from the stack into EDI, push the same value back and then return to where the last pushed value on the stack is pointing to.

This may seem a bit weird, but now EDI contains EIP, and we need that!

So how do we find out the binary values of POP EDI, PUSH EDI, RETN and NOP? Quite simple, we write these in our debugger one by one but don't execute them. This results in these binary opcodes: 5F 57 C3 90

Now don't forget that these values needs to be in Little Endian order, so reverse them till they look like this: 90 C3 57 5F and then implement them in your PUSH opcode, and write it in your debugger like this: PUSH 90C3575F

We also need to create some dummy code to make things a lot easier. The code we're going to use almost looks like all the previous examples we've used, where we use ESI for our null (0) values and EDI for our ASCII pointer which refers to our custom text string.

010FBBB0	68 5F57C390	PUSH 90C3575F
010FBBB5	FFD4	CALL ESP
010FBBB7	90	NOP
010FBBB8	31F6	XOR ESI,ESI
010FBBBA	B8 EA07457E	MOV EAX,7E4507EA
010FBBBF	66:83C7 01	ADD DI,1
010FBBC3	90	NOP
010FBBC4	56	PUSH ESI
010FBBC5	56	PUSH ESI
010FBBC6	57	PUSH EDI
010FBBC7	56	PUSH ESI
010FBBC8	90	NOP
010FBBC9	FFD0	CALL EAX

Figure 5.0.2 – Dummy Popup Code

Now this code does not include our text string yet, but when it's done then we need to place it after the last CALL EAX opcode. I've added a few NOP's in this code to separate the different parts so it is easier to read and understand.

The first 2 opcodes calculates EIP and when the first NOP is hit, EDI contains the value of EIP.

This changes though as we continue to execute our code, so after we've written all these opcodes manually into our debugger we need to adjust "ADD DI, 1" to add the amount of bytes there is from CALL ESP till our text string. In short the space from the first NOP and till the beginning of our text string. This equals 21 in decimal and in hexadecimal, 15.

This is very important to keep in mind that you're working with hexadecimal and not ordinary decimal values. If you forget this like I occasionally do, your shellcode will fail at some point.

Now our text string needs to be encoded into binary (hexadecimal) code before pasting it into our debugger. The string I'm going to use is: "Did you know InterN0T.net is the best community?".

When we've encoded it by e.g. using The XSSOR, removed \x from the string and added a 00-byte it looks like the example below.

44696420796f75206b6e6f7720496e7465724e30542e6e657420697320746865206265737420636f6d6d756e6974793f00

Figure 5.0.3 – Our string encoded in binary form

You may wonder why I did not reverse it into Little Endian? Well we don't need to do that when we're using our text string this way. This saves us some trouble, but there's still the 00-byte at the end which we'll have to deal with later on. But for now let us just see if it works.

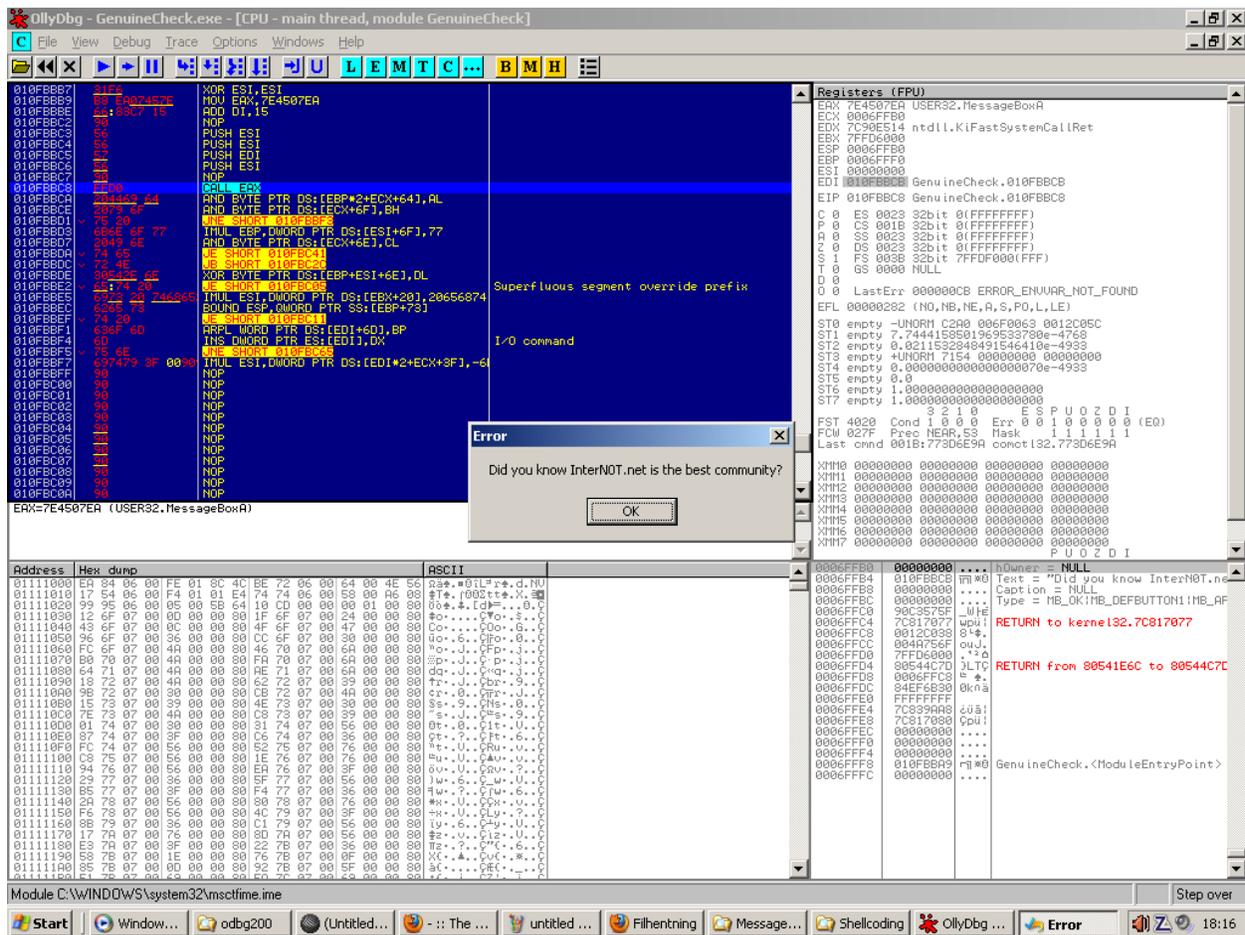


Figure 5.0.4 – Working custom message box

It works as you can see. Don't worry about the "opcodes" after CALL EAX, since we're going to assume that we won't execute these afterwards. If we wanted to make sure this wouldn't happen we could jump somewhere else in memory or perform a system exit call.

There's no need to worry about that, since you can learn this later on if it's required.

Chapter 6

The Easy Way – Part 2

We got our custom message box working in the last chapter but we would like to eliminate the 00-byte which could cause our shellcode to be terminated in a real buffer overflow scenario.

In order to do this we have to either encode that part of our shellcode or encode all of it! This is both fun and interesting to do even though it is also a bit strange at first. We're going to encode a part of our shellcode by "XOR'ing" it, which I have also explained in my other paper about bypassing anti-virus scanners.

In this case we need to use another approach since we have to assume we do not know any hard memory addresses, which means we have to either calculate the start and the end of our string, search for the start and the end and then encode that or what I find the most easy, encode the amount of bytes there is in our shellcode.

First we need to know how big our previous shellcode is:

68	5F	57	C3	90	FF	D4	90	31	F6	B8	EA	07	45	7E	66
83	C7	15	90	56	56	57	56	90	FF	D0	20	44	69	64	20
79	6F	75	20	6B	6E	6F	77	20	49	6E	74	65	72	4E	30
54	2E	6E	65	74	20	69	73	20	74	68	65	20	62	65	73
74	20	63	6F	6D	6D	75	6E	69	74	79	3F	00			

Figure 6.0.1 – Binary Custom Popup Box

The above binary code is exactly 77 bytes long aka 4D in hexadecimal. This is very useful as you'll see when we implement our custom XOR encoder which will also function as a decoder, making things easier for us.

In essence our encoder (and decoder) will have to find out the memory address of the first byte to encode and then loop through X amount of bytes which will be XOR'd. These amounts of bytes are equal to how big our custom popup box is.

In order to find the first memory address of the first byte to XOR encode, we're going to use our previous method used to calculate EIP and then add a value to it and thereby adjusting it.

Now before I go ahead and explain this, a typical problem with XOR encoding in the “section” of the program you’re writing your custom opcodes to, is that it’s not writable by the program.

If it isn’t writable you’ll get an Access Violation Error, and the “XOR” opcode won’t work.

Therefore open the Memory Map in OllyDbg by pressing ALT+M or the “M” icon. Find the name of the executable file you’re injecting your code into, and look for “PE Header”.

Then double-click this and scroll down to where you see “.text” mentioned again.

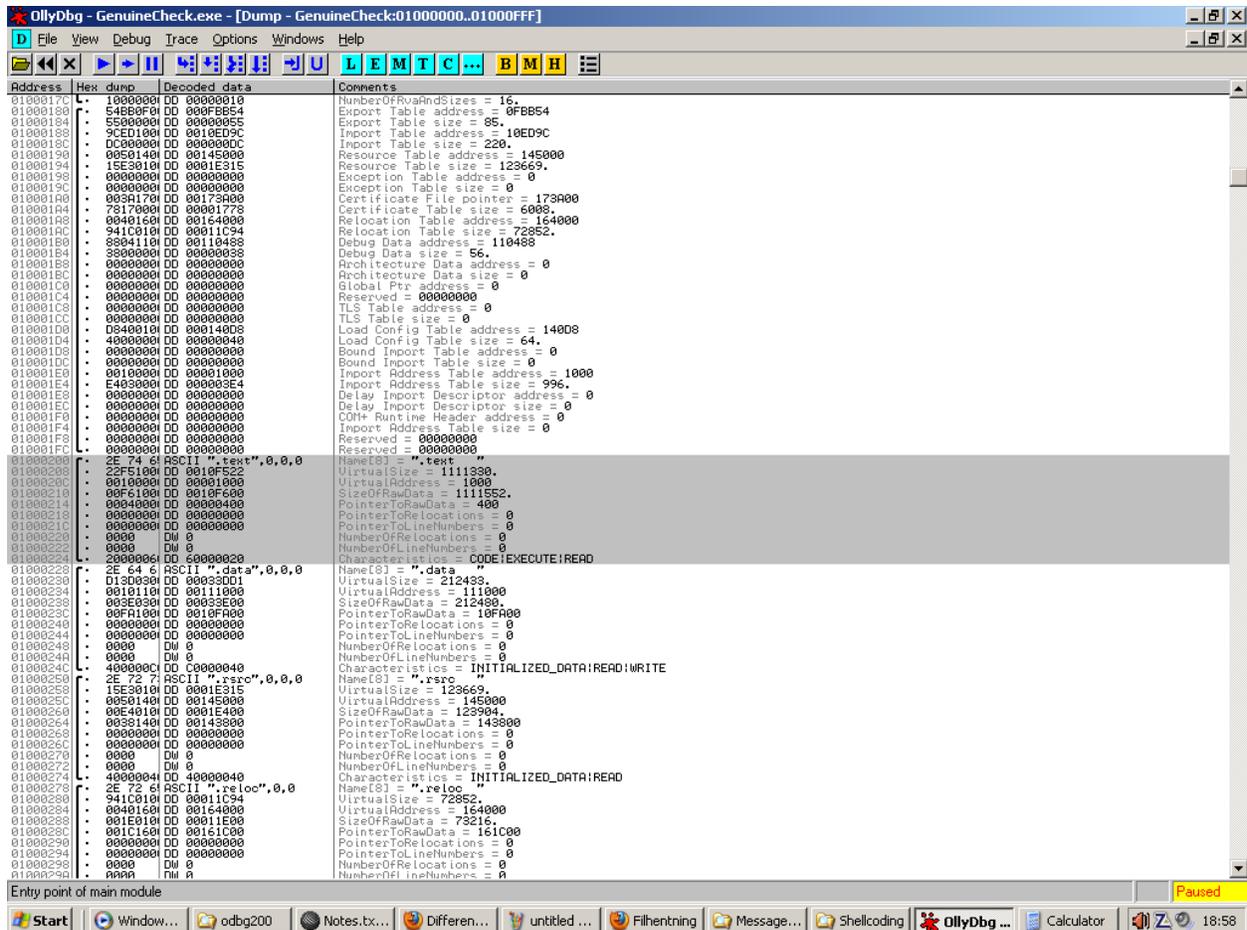


Figure 6.0.2 – PE Header of our executable file

Now double click the “Characteristics” line and edit only the first hexadecimal character to E.

This will make the .text section executable, readable and most importantly: Writeable.

If you want to know more about these sections and how they function, then I suggest you read my paper about bypassing anti-virus scanners but also read up on PE files since this should explain everything you may want to know about this topic.

Now in order to use this change, we need to save these changes and open the new file.

Do this by right clicking the marked and changed line, browse to “Edit” and select: “Copy to Executable”. When you’re in this new window, right click and choose “Save file”. Now open the newly created file which you hopefully saved under a new filename.

After we’ve done this we’re ready to write our custom encoder with the following code.

010FBBA9	68 5F57C390	PUSH 90C3575F
010FBBAE	FFD4	CALL ESP
010FBBB0	66:83C7 10	ADD DI,10
010FBBB4	31C9	XOR ECX,ECX
010FBBB6	80C1 4D	ADD CL,4D
010FBBB9	8037 0D	XOR BYTE PTR DS:[EDI],0D
010FBBBC	47	INC EDI
010FBBBD	E2 FA	LOOP SHORT 010FBBB9

Figure 6.0.3 – Custom XOR Encoder

You might remember the first line, this calculates the EIP and at the third line where EDI contain the current EIP, we add 16! Yes that’s six-teen, not 10 because we’re working with hexadecimal numbers and not decimal numbers. This is very important to keep in mind.

Then ECX is zeroed out (nulled) so we can use this to count down from 77 bytes. ADD CL, 4D adds 77 to ECX and the reason why we use CL and not ECX is to avoid 00 bytes in our shellcode since ADD ECX takes 4 bytes as input while CL takes 1 byte.

XOR BYTE PTR DS:[EDI],0D. Now that may be giving you headache but it’s actually quite simple. Another explanation of this instruction could be: XOR the Byte with 0D, which EDI is pointing to. This will change the byte to another value, which is good to avoid Anti-Virus scanners and also in order to avoid 00 bytes.

When this XOR instruction is executed the first time, EDI is pointing to the byte right after the LOOP opcode. After XOR has been executed, EDI is increased by 1 so it points to the next byte. Then a jump is performed back to “XOR” and 1 is deducted from ECX.

This procedure continues until ECX is equal to 0. When ECX is 0, the “LOOP Jump” is not taken and therefore the rest of the shellcode is executed.

Now this encoder will mess up our popup box completely, but it’ll encode it and avoid the 00-byte which is our goal. When this code has been encoded, we can copy this new and much obfuscated code, and when that is run it automatically decodes itself without any changes made.

It isn’t that efficient in this case, but if our payload containing a lot of 00-bytes it would be very efficient in my opinion since it’s hard to write any smaller encoder than that.

The entire un-encoded shellcode looks like this:

Address	Hex dump	Command
010FBBA9	68 5F57C390	PUSH 90C3575F
010FBBAE	FFD4	CALL ESP
010FBBB0	66:83C7 10	ADD DI,10
010FBBB4	31C9	XOR ECX,ECX
010FBBB6	80C1 4D	ADD CL,4D
010FBBB9	8037 0D	XOR BYTE PTR DS:[EDI],0D
010FBBBC	47	INC EDI
010FBBBD	E2 FA	LOOP SHORT 010FBBB9
010FBBBF	90	NOP
010FBBC0	90	NOP
010FBBC1	68 5F57C390	PUSH 90C3575F
010FBBC6	FFD4	CALL ESP
010FBBC8	90	NOP
010FBBC9	31F6	XOR ESI,ESI
010FBBCB	B8 EA07457E	MOV EAX,7E4507EA
010FBBD0	66:83C7 15	ADD DI,15
010FBBD4	90	NOP
010FBBD5	56	PUSH ESI
010FBBD6	56	PUSH ESI
010FBBD7	57	PUSH EDI
010FBBD8	56	PUSH ESI
010FBBD9	90	NOP
010FBBDA	FFD0	CALL EAX
010FBDBC	204469 64	AND BYTE PTR DS:[EBP*2+ECX+64],AL
010FBBE0	2079 6F	AND BYTE PTR DS:[ECX+6F],BH
010FBBE3	75 20	JNE SHORT 010FBC05
010FBBE5	6B6E 6F 77	IMUL EBP,DWORD PTR DS:[ESI+6F],77
010FBBE9	2049 6E	AND BYTE PTR DS:[ECX+6E],CL
010FBDEC	74 65	JE SHORT 010FBC53
010FBDEE	72 4E	JB SHORT 010FBC3E
010FBBF0	30542E 6E	XOR BYTE PTR DS:[EBP+ESI+6E],DL
010FBBF4	65:74 20	JE SHORT 010FBC17
010FBBF7	6973 20 74686	IMUL ESI,DWORD PTR DS:[EBX+20],20656874
010FBBFE	6265 73	BOUND ESP,QWORD PTR SS:[EBP+73]
010FBC01	74 20	JE SHORT 010FBC23
010FBC03	636F 6D	ARPL WORD PTR DS:[EDI+6D],BP
010FBC06	6D	INS DWORD PTR ES:[EDI],DX
010FBC07	75 6E	JNE SHORT 010FBC77
010FBC09	697479 3F 009	IMUL ESI,DWORD PTR DS:[EDI*2+ECX+3F],-6F

Figure 6.0.8 – Un-Encoded Shellcode

While the encoded shellcode looks like this:

Address	Hex dump	Command
010FBBA9	68 5F57C390	PUSH 90C3575F
010FBBAE	FFD4	CALL ESP
010FBBB0	66:83C7 10	ADD DI,10
010FBBB4	31C9	XOR ECX,ECX
010FBBB6	80C1 4E	ADD CL,4E
010FBBB9	8037 0D	XOR BYTE PTR DS:[EDI],0D
010FBBBC	47	INC EDI
010FBBBD	E2 FA	LOOP SHORT 010FBBB9
010FBBBF	90	NOP
010FBBC0	9D	POPF
010FBBC1	65:52	PUSH EDX
010FBBC3	5A	POP EDX
010FBBC4	CE	INTO
010FBBC5	9D	POPF
010FBBC6	F2:D99D 3CFBB	REPNE FSTP DWORD PTR SS:[EBP+E7B5FB3C]
010FBBCD	0A48 73	OR CL,BYTE PTR DS:[EAX+73]
010FBBD0	6B8E CA189D5B	IMUL ECX,DWORD PTR DS:[ESI+5B9D18CA],5B
010FBBD7	5A	POP EDX
010FBBD8	5B	POP EBX
010FBBD9	9D	POPF
010FBBDA	F2	REPNE
010FBBDB	DD	DB DD
010FBBDC	2D 4964692D	SUB EAX,2D696449
010FBBE1	74 62	JE SHORT 010FBC45
010FBBE3	78 2D	JS SHORT 010FBC12
010FBBE5	66:6362 7A	ARPL WORD PTR DS:[EDX+7A],SP
010FBBE9	2D 44637968	SUB EAX,68796344
010FBBEE	7F 43	JG SHORT 010FBC33
010FBBF0	3D 59236368	CMP EAX,68632359
010FBBF5	79 2D	JNS SHORT 010FBC24
010FBBF7	64:7E 2D	JLE SHORT 010FBC27
010FBBFA	79 65	JNS SHORT 010FBC61
010FBBFC	68 2D6F687E	PUSH 7E686F2D
010FBC01	79 2D	JNS SHORT 010FBC30
010FBC03	6E	OUTS DX,BYTE PTR DS:[ESI]
010FBC04	6260 60	BOUND ESP,QWORD PTR DS:[EAX+60]
010FBC07	78 63	JS SHORT 010FBC6C
010FBC09	64:79 74	JNS SHORT 010FBC80
010FBC0C	320D 90909090	XOR CL,BYTE PTR DS:[90909090]

Figure 6.0.9 – Encoded Shellcode

Ending Words

Knowing Assembly in order to write and create shellcode directly is indeed a very good idea since it allows the ethical hacker to create efficient, minimalistic and optimized shellcode for future exploitation scenarios which may require hardcore expertise.

Therefore you should if you don't already know Assembly, want to learn more. Don't go hardcore but try to create your own shellcode which calls a socket and listens for connections, or perhaps execute a system command (calc.exe?), or maybe another function from the API.

If you play long enough with the language it won't be as confusing, most of the time.

References

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